

# Mini-course on Stationary Iterative Methods

Lecture 2 - Fixed point methods

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**∷** Outline



# Definitions and basic notions

# Iterative methods for linear problems

# Definitions and basic notions

# **O** Iterative methods



Suppose we want to solve the nonlinear system

$$F(\mathbf{x}) = \begin{bmatrix} f_1(x_1, x_2 \dots, x_n) \\ f_2(x_1, x_2 \dots, x_n) \\ \vdots \\ f_n(x_1, x_2 \dots, x_n) \end{bmatrix} = 0 \quad \text{for} \quad \mathbf{x} = \begin{bmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{bmatrix}.$$

We assume that  $F: D \subset \mathbb{R}^n \to \mathbb{R}^n$  is (Fréchet-) differentiable.

- If  $F(\mathbf{x}) = A\mathbf{x} \mathbf{b}$ , then the system is linear and can be solved by e.g. Gaussian elimination. Iterative methods can also be used, see the next section.
- If F is nonlinear, there is no *direct* method that solves  $F(\mathbf{x}) = 0$  in a finite number of operations for *arbitrary* F.
- An iterative method generates a sequence of approximations  $\mathbf{x}^1, \mathbf{x}^2, \ldots$ , such that  $\mathbf{x}^k$  (hopefully) converges to the solution  $\mathbf{x}^*$  as  $k \to \infty$ .



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$$F(\mathbf{x}) \approx F(\mathbf{x}^0) + F'(\mathbf{x}^0)(\mathbf{x} - \mathbf{x}^0) + O(\|\mathbf{x} - \mathbf{x}^0\|^2).$$

Here,  $F'(\mathbf{x}^0)$  is the Jacobian matrix :

$$F'(\mathbf{x}_0) = \begin{bmatrix} \frac{\partial f_1}{\partial x_1}(\mathbf{x}^0) & \frac{\partial f_1}{\partial x_2}(\mathbf{x}^0) & \cdots & \frac{\partial f_1}{\partial x_n}(\mathbf{x}^0) \\ \frac{\partial f_2}{\partial x_1}(\mathbf{x}^0) & \frac{\partial f_2}{\partial x_2}(\mathbf{x}^0) & \cdots & \frac{\partial f_2}{\partial x_n}(\mathbf{x}^0) \\ \vdots & \vdots & \vdots \\ \frac{\partial f_n}{\partial x_1}(\mathbf{x}^0) & \frac{\partial f_n}{\partial x_2}(\mathbf{x}^0) & \cdots & \frac{\partial f_n}{\partial x_n}(\mathbf{x}^0) \end{bmatrix}$$



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2. Find the root of the *linearized* equation and call it  $\mathbf{x}^1$  :

$$0 = F(\mathbf{x}^0) + F'(\mathbf{x}^0)(\mathbf{x}^1 - \mathbf{x}^0) \implies \mathbf{x}^1 = \mathbf{x}^0 - (F'(\mathbf{x}^0))^{-1}F(\mathbf{x}^0).$$



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3. Repeat the process using the newly calculated point :

$$\mathbf{x}^{k+1} = \mathbf{x}^k - (F'(\mathbf{x}^k))^{-1}F(\mathbf{x}^k), \qquad k = 1, 2, \dots$$



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4. Stop when  $\mathbf{x}^k$  and  $\mathbf{x}^{k+1}$  are very close (or when  $||F(\mathbf{x}^k)||$  is very small, or both).



Apply Newton's method to the real and imaginary parts of the equation  $z^3 - 1 = 0$  for z = x + iy:

$$F(x,y) = \begin{bmatrix} f_1(x,y) \\ f_2(x,y) \end{bmatrix} = \begin{bmatrix} x^3 - 3xy^2 - 1 \\ 3x^2y - y^3 \end{bmatrix} = 0$$

500 600

700 -800 -900 -



Newton attraction basins



# Newton's method – remarks



- Instead of computing the matrix inverse in  $\mathbf{x}^{k+1} = \mathbf{x}^k (F'(\mathbf{x}^k))^{-1}F(\mathbf{x}^k)$  the Newton step is usually implemented as follows :
  - 1. Solve  $F'(\mathbf{x}^k)\mathbf{z} = F(\mathbf{x}^k)$  for  $\mathbf{z}$  using an appropriate linear solver;

2. Update 
$$\mathbf{x}^{k+1} = \mathbf{x}^k - \mathbf{z}$$
.

A good linear solver for F' is therefore crucial !

- Newton's method can diverge or stagnate when  $\mathbf{x}_0$  is too far from the solution !
- Jacobian evaluation is expensive; one could freeze the Jacobian at the initial guess (or update every few iterations) to obtain

$$\mathbf{x}^{k+1} = \mathbf{x}^k - (F'(\mathbf{x}^0))^{-1}F(\mathbf{x}^k), \qquad k = 0, 1, 2, \dots$$

This is known as the simplified Newton method.

# • Fixed point methods



• Newton's method can be written in the general form  $\mathbf{x}^{k+1} = G(\mathbf{x}^k)$ , where  $G: D \subset \mathbb{R}^n \to \mathbb{R}^n$  is the **fixed point map**. For Newton, we have

$$G(\mathbf{x}) = \mathbf{x} - (F'(\mathbf{x}))^{-1}F(\mathbf{x}).$$

• If a fixed point method converges and G is continuous, then the limit point  $\mathbf{x}^*$  satisfies

$$\mathbf{x}^* = \lim_{k \to \infty} \mathbf{x}^k = \lim_{k \to \infty} G(\mathbf{x}^{k-1}) = G\left(\lim_{k \to \infty} \mathbf{x}^{k-1}\right) = G(x^*).$$

The method therefore converges to a **fixed point** of G.

• Example : the fixed point for Newton's method is

$$\mathbf{x}^* = G(\mathbf{x}^*) = \mathbf{x}^* - (F'(\mathbf{x}^*))^{-1} \mathcal{F}(\mathbf{x}^*) \implies (F'(\mathbf{x}^*))^{-1} F(\mathbf{x}^*) = 0.$$

If  $F'(\mathbf{x}^*)$  is non-singular, then  $F(\mathbf{x}^*) = 0 \implies \mathbf{x}^*$  is a solution to the nonlinear system.

#### • Convergence of fixed point methods



**Definition :** Let  $\mathbf{x}^*$  be a fixed point of G and let  $\mathbf{x}^0, \mathbf{x}^1, \ldots$  be iterates that satisfy  $\mathbf{x}^{k+1} = G(\mathbf{x}^k)$ ,  $k = 0, 1, 2, \ldots$  Then the **error** of the kth iterate is defined as

$$\mathbf{e}^k := \mathbf{x}^k - \mathbf{x}^*.$$

To understand how the error behaves as a function of k, we calculate

$$\begin{aligned} \mathbf{e}^{k+1} &= \mathbf{x}^{k+1} - \mathbf{x}^* \\ &= G(\mathbf{x}^k) - G(\mathbf{x}^*) \\ &= G(\mathbf{x}^*) + G'(\mathbf{x}^*)(\mathbf{x}^k - \mathbf{x}^*) + O(\|\mathbf{x}^k - \mathbf{x}^*\|^2) - G(\mathbf{x}^*) \\ &= G'(\mathbf{x}^*)\mathbf{e}^k + O(\|\mathbf{e}^k\|^2). \end{aligned}$$

# • Example 1



To find one of the roots of  $x^3 - 3x + 1 = 0$ , we start with  $x_0 = 0$  and consider the method

$$x^{k+1} = \frac{1}{3}((x^k)^3 + 1).$$

Then  $G(x) = \frac{1}{3}(x^3+1)$ ,  $G'(x) = x^2$ . For the fixed point  $x^* \approx 0.3473$ , we have  $G'(x^*) \approx 0.1206$ , so for  $x^k$  close to  $x^*$ , we have

 $e^{k+1} \approx 0.1206 \, e^k + O((e^k)^2),$ 

so the error is reduced by about a factor of  $\boldsymbol{8}$  with each additional iteration.



k	$x^k$	$e^k$	$e^{k+1}/e^k$
0	0	-0.347296355333861	0.040204919476061
1	0. <mark>3</mark> 33333333333333333333333333333333333	-0.013963022000527	0.115830440439082
2	0. <mark>34</mark> 5679012345679	-0.001617342988182	0.120053933035803
3	0. <mark>347</mark> 102186947062	-0.000194168386799	0.120547337021927
4	0. <mark>3472</mark> 72948851898	-0.000023406481962	0.120606629621737
5	0. <mark>34729</mark> 3532356960	-0.000002822976901	0.120613778000051
6	0. <mark>347296</mark> 014843951	-0.000000340489909	0.120614640035547
7	0. <mark>3472963</mark> 14265793	-0.000000041068068	0.120614743217285
8	0. <mark>34729635</mark> 0380446	-0.000000004953414	0.120614753443415
9	0. <mark>34729635</mark> 4736406	-0.000000000597455	0.120614636167759
10	0. <mark>347296355</mark> 261799	-0.00000000072062	-

### • What about Newton's method?



In one dimension :

$$G(x) = x - \frac{f(x)}{f'(x)}$$
  

$$G'(x) = 1 - \frac{f'(x) \cdot f'(x) - f(x)f''(x)}{(f'(x))^2} = \frac{f(x)f''(x)}{(f'(x))^2}$$

• If  $x^*$  is a root of f(x) such that  $f'(x^*) \neq 0$  (i.e., a simple root), then

$$G'(x^*) = \frac{0 \cdot f''(x^*)}{(f'(x^*))^2} = 0.$$

So for  $|x^k - x^*|$  small, we have

$$e^{k+1} = G'(x^*)e^k + O(|e^k|^2) = O(|e^k|^2)$$

Thus, we have quadratic convergence !



To solve  $x^3 - 3x + 1 = 0$  using Newton's method, we write

$$x^{k+1} = x^k - \frac{f(x^k)}{f'(x^k)} = x^k - \frac{(x^k)^3 - 3x^k + 1}{3(x^k)^2 - 3}$$

The number of correct digits **doubles** after every iteration  $\implies$  quadratic convergence

k	$x^k$	$e^k$	$e^{k+1}/(e^k)^2$
0	0	-0.347296355333861	-0.115765451777953
1	0. <mark>3</mark> 33333333333333333333333333333333333	-0.013963022000527	-0.380236133250081
2	0. <mark>3472</mark> 222222222222	-0.000074133111638	-0.394851507080994
3	0. <mark>34729635</mark> 3163868	-0.000000002169993	0
4	0.347296355333861	0	_

# Some problems with Newton



• Recall :

$$e^{k+1} = G'(x^*)e^k + O(|e^k|^2) = O(|e^k|^2)$$

If  $|e^k|$  is large, then

$$|e^{k+1}| \approx C |e^k|^2 > |e^k| \qquad \text{if } |Ce^k| > 1.$$

So Newton's method may not converge in this case; it may cycle or diverge.

• If  $x^*$  is a **double root** of f, i.e., if  $f(x^*) = f'(x^*) = 0$  (but  $f''(x^*) \neq 0$ ), then

$$\lim_{x \to x^*} G'(x) = \lim_{x \to x^*} f''(x) \frac{f(x)}{(f'(x))^2} = f''(x^*) \lim_{x \to x^*} \frac{f'(x)}{2f'(x)f''(x)} = \frac{1}{2}.$$

So Newton converges **linearly** with 1/2 as the asymptotic contraction factor.

• A more involved analysis leads to analogous results for higher dimensions

#### **O** Problems with Newton : some fixes



- Line search :
  - Instead of using the full Newton correction  $\mathbf{x}^{k+1} = \mathbf{x}_k (F'(\mathbf{x}^k))^{-1}F(\mathbf{x}^k)$ , make a shorter step

$$\mathbf{x}^{k+1} = \mathbf{x}_k - \boldsymbol{\alpha}(F'(\mathbf{x}^k))^{-1}F(\mathbf{x}^k), \qquad 0 < \alpha < 1$$

- $\alpha$  is chosen to ensure  $\|F(\mathbf{x}^{k+1})\| < \|F(\mathbf{x}^k)\|$
- Under some hypotheses, one can show convergence to a solution  $F(\mathbf{x}^*) = 0$ .
- Preconditioning : transform F(x) = 0 to some equivalent problem H(x) = 0 with the same solution  $\mathbf{x}^*$ , but which is easier to solve for Newton (e.g., closer to a linear function)
- Continuation : If the problem depends continuously on a parameter  $\alpha,$  i.e.,

$$F(x(\alpha);\alpha) = 0,$$

start with an "easy"  $\alpha$  and vary  $\alpha$  step by step, using the solution from the previous step as an initial guess

Iterative methods for linear problems

#### Suppose we want to solve the linear system

$$\begin{bmatrix} 2 & -1 & & 0 \\ -1 & 2 & -1 & & \\ & \ddots & \ddots & \ddots & \\ & & -1 & 2 & -1 \\ 0 & & & -1 & 2 \end{bmatrix} \mathbf{x} = \begin{bmatrix} f_1 \\ f_2 \\ \vdots \\ \vdots \\ f_n \end{bmatrix}.$$

Instead of eliminating unknowns, Jacobi (1845) has the following idea :

- 1. Make some initial guess  $\mathbf{x}^0 = (x_1^0, x_2^0, \dots, x_n^0)^T$
- 2. Use the *i*th equation to solve for the new  $x_i$  using the previous data :

$$-x_{i-1} + 2x_i - x_{i+1} = f_i \implies x_i^{k+1} = \frac{1}{2}(x_{i-1}^k + x_{i+1}^k + f_i), \qquad k = 1, \dots, n$$

3. Iterate until convergence (?!).





































# ● Gauss-Seidel method



Make a programming mistake :

• Jacobi :

$$x_i^{k+1} = \frac{1}{2}(x_{i-1}^k + x_{i+1}^k + f_i), \qquad i = 1, \dots, n$$

• Gauss-Seidel :

$$x_i^{k+1} = \frac{1}{2}(x_{i-1}^{k+1} + x_{i+1}^k + f_i), \qquad i = 1, \dots, n$$

Gauss explained in a letter in 1823 how to use this method to solve a least squares problem :

"I recommend this method to you for imitation. You will hardly ever again eliminate directly, at least not when you have more than 2 unknowns. The indirect procedure can be done while half asleep, or while thinking about other things."

**\bigcirc** Gauss-Seidel method, n = 10







**\bigcirc** Gauss-Seidel method, n = 10

























#### Stationary iterative method : General form



The Jacobi and Gauss-Seidel methods for solving  $A\mathbf{x} = \mathbf{f}$  can be rewritten in the general form

$$M\mathbf{x}^{k+1} = N\mathbf{x}^k + \mathbf{f},$$

where M is a non-singular matrix, and N = M - A.

• Jacobi : 
$$M = D = \text{diag}(a_{11}, a_{22}, \dots, a_{nn})$$

• Gauss-Seidel : 
$$M = L + D = \begin{bmatrix} a_{11} & & \\ a_{21} & a_{22} & \\ \vdots & \ddots & \\ a_{n1} & \cdots & \cdots & a_{nn} \end{bmatrix}$$
, i.e., the lower triangular part of  $A$ 

• Other choices are possible, as long as systems involving M are cheap to solve. Note that we need N = M - A so that the fixed point  $\mathbf{x}^* = \lim \mathbf{x}^k$  satisfies

$$M\mathbf{x}^* = N\mathbf{x}^* + \mathbf{f} \iff (M - N)\mathbf{x}^* = \mathbf{f} \iff A\mathbf{x}^* = \mathbf{f}.$$

$$M\mathbf{x}^{k+1} = N\mathbf{x}^k + \mathbf{f}$$

• This is a fixed point method, because we can write it in terms of a fixed point map  $\mathcal{G}(\mathbf{x})$  :

$$\mathbf{x}^{k+1} = M^{-1}N\mathbf{x}^k + M^{-1}\mathbf{f} =: \mathcal{G}(\mathbf{x}^k),$$

• Another way to write this is in terms of the **residual**  $\mathbf{r}^k = \mathbf{f} - A\mathbf{x}^k$  :

$$M\mathbf{x}^{k+1} = (M-A)\mathbf{x}^k + \mathbf{f} = M\mathbf{x}^k + \mathbf{r}^k \implies \mathbf{x}^{k+1} = \mathbf{x}^k + M^{-1}\mathbf{r}^k,$$

so the method will keep updating  $\mathbf{x}^{k+1}$  as long as  $\mathbf{r}^k \neq 0$ .

• A stopping criterion is usually needed to decide when the solution  $\mathbf{x}^k$  is "close enough". Some options are  $\|\mathbf{r}^k\|$ ,  $\|\mathbf{x}^{k+1} - \mathbf{x}^k\|$ , or some other norm of the error/residual.

# **O** Direct vs iterative method



- $\bullet~\mbox{Recall}$  : cost of direct methods depends heavily on the density of LU factors
- Iterative methods can be matrix free, with much lower cost per iteration
- But : speed of convergence is paramount !



# • Convergence of a stationary iterative method



$$M\mathbf{x}^{k+1} = N\mathbf{x}^k + \mathbf{f} \tag{(*)}$$

- Define the error  $\mathbf{e}^k = \mathbf{x}^k \mathbf{x}^*$ .
- Subtracting the above from (\*) gives

$$M(\mathbf{x}^{k+1} - \mathbf{x}^*) = N(\mathbf{x}^k - \mathbf{x}^*) \implies M\mathbf{e}^{k+1} = N\mathbf{e}^k.$$

• Therefore, we have

$$e^{k+1} = Ge^k$$
, where  $G = M^{-1}N = I - M^{-1}A$ .

• Applying the above relation several times yields

$$\mathbf{e}^{k} = G\mathbf{e}^{k-1} = G(G\mathbf{e}^{k-2}) = G^{2}(G\mathbf{e}^{k-3}) = \dots = G^{k}\mathbf{e}^{0},$$

where  $e^0 = x^0 - x^*$  is the initial error (generally unknown).

Let  $e^0 \neq 0$  be an **eigenvector** of G, i.e.,  $Ge^0 = \lambda e^0$  for some  $\lambda$  (which may be complex). Then

$$\mathbf{e}^{k} = G^{k} \mathbf{e}^{0} = \lambda G^{k-1} \mathbf{e}^{0} = \dots = \lambda^{k} \mathbf{e}^{0} \implies \|\mathbf{e}^{k}\| = |\lambda|^{k} \|\mathbf{e}^{0}\|$$

Several cases :

- If  $|\lambda| < 1$ , then  $|\lambda|^k \to 0$  as  $k \to \infty \implies ||\mathbf{e}^k|| \to 0$ , i.e., convergence !
- If  $|\lambda| > 1$ , then  $|\lambda|^k \to \infty \implies \|\mathbf{e}^k\| \to \infty$ , i.e., divergence.
- If  $|\lambda| = 1$ , then  $\|\mathbf{e}^k\| = \|\mathbf{e}^0\|$  for all  $k \implies$  no convergence to  $\mathbf{x}^*$ .

In general,  $e^0$  is a (usually unknown) linear combination of eigenvectors of G.

**Theorem :** A stationary iterative method  $M\mathbf{x}^{k+1} = N\mathbf{x}^k + \mathbf{f}$  converges to the unique solution  $\mathbf{x}^* = \mathbf{A}^{-1}\mathbf{f}$  for all initial guess  $\mathbf{x}^0$  if and only if **all** the eigenvalues of  $G = M^{-1}N$  have modulus strictly less than 1.

# Onvergence rate



 ${\rm Definition}: {\rm The \ spectral \ radius}\ \rho(G)$  of a matrix G is defined as the maximum modulus of all the eigenvalues, i.e.,

$$p(G) = \max_{1 \le i \le n} |\lambda_i|,$$

where  $\lambda_1, \ldots, \lambda_n$  are the eigenvalues of G.

Let  $\rho(G) < 1$ , so that the stationary method converges. If G is diagonalizable with eigenvectors  $\{\mathbf{v}_1, \mathbf{v}_2, \dots, \mathbf{v}_n\}$ , then we have

$$\mathbf{e}^{0} = \sum_{i=1}^{n} c_{i} \mathbf{v}_{i} \implies \mathbf{e}^{k} = G^{k} \mathbf{e}^{0} = \sum_{i=1}^{n} c_{i} \lambda_{i}^{k} \mathbf{v}_{i}$$
$$|\mathbf{e}^{k}|| \leq \sum_{i=1}^{n} |c_{i}| |\lambda_{i}|^{k} ||\mathbf{v}_{i}|| \leq C(\rho(G))^{k}$$

Therefore, the error norm is multiplied by roughly  $\rho(G)$  at each iteration.

$$\|\mathbf{e}^k\| \leq C(\rho(G))^k$$

- $\rho(G)$  gives the contraction factor of the method :
  - $\bullet \ \rho(G) \ll 1 \implies {\rm fast \ convergence}$
  - $\rho(G) \approx 1 \implies$  slow convergence
- The "best" method is when  $\rho(G) = 0$ , which happens when G = 0:

$$G = I - M^{-1}A = 0 \iff M^{-1}A = I \iff M = A.$$

But solving with M would be just as expensive as solving with A !

- Designing a good method involves a tradeoff :
  - *M* is a "good" approximation of *A* (in some sense)
  - Systems involving M are cheap(er) to solve than those involving A

# Sexample : Jacobi and Gauss-Seidel



Consider the  $10\times10$  system

$$\begin{bmatrix} a & -1 & & \\ -1 & a & -1 & & \\ & \ddots & \ddots & \ddots & \\ & & -1 & a & -1 \\ & & & -1 & a \end{bmatrix} \mathbf{x} = \begin{bmatrix} 1 \\ 1 \\ \vdots \\ \vdots \\ 1 \end{bmatrix}$$

for a = 2 and a = 4.

**Example : Jacobi,** a = 2





A tenfold reduction in error requires  $\rho^k = 0.1 \implies k = 55.68$  iterations !

**\bigcirc** Example : Jacobi, a = 4





A tenfold reduction in error requires  $\rho^k = 0.1 \implies k = 3.13$  iterations !

**Solution** Example : Gauss-Seidel, a = 2





A tenfold reduction in error requires  $\rho^k = 0.1 \implies k = 27.84$  iterations !

**\bigcirc** Example : Gauss-Seidel, a = 4





A tenfold reduction in error requires  $\rho^k = 0.1 \implies k = 1.57$  iterations!



**Theorem :** Let A be strictly diagonally dominant, i.e., the diagonal element in each row is greater than the sum of the off-diagonal entries (in absolute value) in that row. Then the Jacobi and Gauss-Seidel methods, applied to  $A\mathbf{x} = \mathbf{f}$ , both converge for all initial guess.

**Theorem :** Let A be symmetric positive definite. Then the Gauss-Seidel method, applied to  $A\mathbf{x} = \mathbf{f}$ , converges for all initial guess.

**Note :** for the matrices in our example, it can be shown that the eigenvalues of the Gauss-Seidel matrix are the **squares** of the eigenvalues of the Jacobi matrix. Thus, Gauss-Seidel converges twice as fast as Jacobi in those examples.

# Block methods



Suppose we have a coupled multiphysics problem :

$$\begin{cases} F(\mathbf{u}, \mathbf{v}) = 0, \\ G(\mathbf{u}, \mathbf{v}) = 0. \end{cases}$$

If we solve this problem by Newton's method, we need to solve linear systems of the form

$$\underbrace{\begin{bmatrix} F_{\mathbf{u}}(\mathbf{u}^{j},\mathbf{v}^{j}) & F_{\mathbf{v}}(\mathbf{u}^{j},\mathbf{v}^{j}) \\ G_{\mathbf{u}}(\mathbf{u}^{j},\mathbf{v}^{j}) & G_{\mathbf{v}}(\mathbf{u}^{j},\mathbf{v}^{j}) \end{bmatrix}}_{A} \underbrace{\begin{bmatrix} \mathbf{u}^{j+1} - \mathbf{u}^{j} \\ \mathbf{v}^{j+1} - \mathbf{v}^{j} \end{bmatrix}}_{\mathbf{x}} = \underbrace{- \begin{bmatrix} F(\mathbf{u}^{j},\mathbf{v}^{j}) \\ G(\mathbf{u}^{j},\mathbf{v}^{j}) \end{bmatrix}}_{\mathbf{f}}.$$

Block Jacobi is defined as  $M\mathbf{x}^{k+1} = N\mathbf{x}^k + \mathbf{f}$ , where M contains only a single "physics" :

$$M := \begin{bmatrix} F_{\mathbf{u}}(\mathbf{u}^j, \mathbf{v}^j) & 0\\ 0 & G_{\mathbf{v}}(\mathbf{u}^j, \mathbf{v}^j) \end{bmatrix}.$$

In other words, within each Newton step j, we solve for k = 0, 1, 2, ...

$$\begin{split} F_{\mathbf{u}}(\mathbf{u}^{j},\mathbf{v}^{j})(\mathbf{u}^{j+1,k+1}-\mathbf{u}^{j}) &= -F(\mathbf{u}^{j},\mathbf{v}^{j}) - F_{\mathbf{v}}(\mathbf{u}^{j},\mathbf{v}^{j})(\mathbf{v}^{j+1,k}-\mathbf{v}^{j}),\\ G_{\mathbf{v}}(\mathbf{u}^{j},\mathbf{v}^{j})(\mathbf{v}^{j+1,k+1}-\mathbf{v}^{j}) &= -G(\mathbf{u}^{j},\mathbf{v}^{j}) - G_{\mathbf{u}}(\mathbf{u}^{j},\mathbf{v}^{j})(\mathbf{u}^{j+1,k}-\mathbf{u}^{j}). \end{split}$$

One can also use block Gauss-Seidel to maintain some coupling by keeping one off-diagonal block :

$$M := \begin{bmatrix} F_{\mathbf{u}}(\mathbf{u}, \mathbf{v}) & 0\\ G_{\mathbf{u}}(\mathbf{u}, \mathbf{v}) & G_{\mathbf{v}}(\mathbf{u}, \mathbf{v}) \end{bmatrix}.$$

Of course, one could also consider a **nonlinear** block Jacobi or Gauss-Seidel fixed-point iteration by solving directly the nonlinear problems

$$F(\mathbf{u}^{k+1}, \mathbf{v}^k) = 0,$$
  
 $G(\mathbf{u}^k, \mathbf{v}^{k+1}) = 0,$   $k = 0, 1, 2, ...$ 

## **Example : CPR method** (Wallace et al., 1985)



To solve the Jacobian systems arising from two-phase flow equations :

$$\begin{split} \text{Water}: & \quad \frac{\partial}{\partial t}(\phi\rho_wS) = \nabla \cdot (K(\mathbf{x})\rho_w\lambda_w(S)\,\nabla p) + q_w, \\ \text{Oil}: & \quad \frac{\partial}{\partial t}(\phi\rho_o(1-S)) = \nabla \cdot (K(\mathbf{x})\rho_o\lambda_o(S)\,\nabla p) + q_o, \end{split}$$

1. Add the two equations to obtain the discrete pressure equation

$$-\nabla \cdot (K_T(\mathbf{x}, S)\nabla p) = \tilde{q},$$

- 2. Use e.g. algebraic multigrid to solve the pressure equation.
- 3. Compute the new velocity field  $\mathbf{v}_T$ .
- 4. Solve local,  $2 \times 2$  systems to update the saturation S, in decreasing order of pressure, or using an incomplete LU factorization.
- 5. Iterate to convergence.

# • Example : cooling by fluids



Momentum :	$\nabla \cdot (2\eta \nabla^s \mathbf{v}) - \rho_f(\mathbf{v} \cdot \nabla) \mathbf{v} - \nabla p = 0$	in $\Omega_f$ ,
Mass :	$ abla \cdot ( ho_f \mathbf{v}) = 0$	in $\Omega_f$ ,
Heat :	$-\nabla \cdot (K_f \nabla T) + \nabla \cdot (\mathbf{v} \rho_f C_f T) = 0$	in $\Omega_f$ ,
Heat :	$-\nabla \cdot (K_s \nabla T) = q$	in $\Omega_s$ ,
Coupling :	$K_s \nabla T \cdot \mathbf{n}_s + K_f \nabla T \cdot \mathbf{n}_f = 0$	on $\Gamma$ .

Apply block Jacobi/Gauss-Seidel to sub-problems :

- Solve Navier-Stokes for  $(\mathbf{v},p)$  at fixed temperature, and heat equations at fixed  $\mathbf{v}$
- Further split Navier-Stokes into a momentum equation and a pressure equation
- Split the heat equation into two subproblems based on the material

Performance depends on "coupling strength", which depends on  $\rho$ ,  $\mathbf{v}$ , ...

# Implementation

- In the exercise session on Tuesday, you will simulate the cooling problem in MEF++, the finite element software developed by GIREF at Université Laval.
- MEF++ uses  $PETSc^1$  to handle solution to linear problems arising from Newton's method
- PETSc offers block preconditioners under the FieldSplit option. Preconditioners available :
  - Additive Schwarz (Block Jacobi + overlap)
  - Multiplicative Schwarz (Block Gauss-Seidel + overlap)
  - Symmetric Multiplicative Schwarz (Forward + backward Gauss-Seidel)
  - Schur complement preconditioning (Eliminate (1,1)block exactly)
- Try different configurations to see what works best for your problem !



<sup>&</sup>lt;sup>1</sup> https://www.mcs.anl.gov/petsc/